#### Advanced Mathematical Structures

Instructor: Girish Varma • Course Code: IMA410

#### Solve 2 problems • 5 marks each

## 1 Averaging Argument

- Suppose there are *N* members in a library containing *B* books. If every person likes at least 1/3 of the books in the library, show that the library has a book, which at least 1/3 of people like.
   (3 marks)
- 2. Let f: {0,1}<sup>n</sup> → {0,1} be some function. If we have an algorithm A which take inputs x ∈ {0,1}<sup>n</sup>, y ∈ {0,1}<sup>m</sup> such that A(x,y) = f(x) with probability at least ρ where x is chosen at random and y is chosen randomly, independently from {0,1}<sup>m</sup>. Show that there exists a single string y\* ∈ {0,1}<sup>m</sup> such that Pr<sub>x</sub>[A(x, y\*) = f(x)] ≥ ρ (ie. only x is chosen randomly here). (2 marks)

### 2 Hypergraphs and Rainbow Colorings

A *k*-uniform hypergraph is a generalization of a graph  $\mathcal{G}([n], \mathcal{E})$ , which has a vertex set [n] and a hyperedge set  $\mathcal{E} = \{e_1, \dots, e_m\}$  where each  $e_i \subseteq [n]$  of size *k*. Note that for k = 2, we get an undirected graph. An assignment of colors to the vertex set is said to (t, c)-rainbow color an edge *e*, if it uses only *c* distinct colors in total and the number of distinct colors in *e* is at least *t*. A (t, c)-rainbow coloring of hypergraph  $\mathcal{G}([n], \mathcal{E})$  is an assignment that is a (t, c)-rainbow coloring of every  $e \in E$ .

- 1. Show that for a random assignment of colors in [k] to the vertices, the expected number of (k, k)-rainbow colored edges is  $\frac{|\mathcal{E}| \cdot k!}{k^k}$  (2 marks)
- 2. Show that any hypergraph with number of edges  $m < \frac{k^k}{k^k k!}$ , has a (k, k)-rainbow coloring. (1 mark)
- 3. Use the method of conditional expectation, to give a deterministic algorithm to find the (k, k)-rainbow coloring for  $\frac{|\mathcal{E}| \cdot k!}{k^k}$  from first part. (2 marks)

(You can substitute k = 2 and solve all the problems to get 2.5 marks).

# 3 Randomized Perfect Matching

Consider a bipartite graph G(V, E) with vertex set V = [n]. Let A be the adjacency matrix. Let

$$\operatorname{perm}(A) = \sum_{\sigma \in S_n} \left( \prod_{i \in [n]} a_{i\sigma(i)} \right) \quad \text{and} \quad \det(A) = \sum_{\sigma \in S_n} (-1)^{\operatorname{parity}(\sigma)} \left( \prod_{i \in [n]} a_{i\sigma(i)} \right)$$

- 1. Show that perm(A) is equal to the number of perfect matchings in *G*. (1 mark)
- 2. Let *B* be the matrix obtained from *A* by replacing  $A_{ij}$  with  $x_{ij}A_{ij}$  where  $x_{ij}$  are some variables. Note that det(*B*), perm(*B*) are polynomials in these variables. Show that det(*B*)  $\equiv$  0 if and only if there are no perfect matching. (1 marks)
- 3. Let B' be the random matrix obtained by substituting each x<sub>ij</sub> with uniformly and independently chosen values from [2n].
  Show that Pr[det(B') = 0|B ≠ 0] ≤ 1/2.
  (2 marks)
- 4. Using the above give a randomized algorithm with one-sided error for checking whether a graph has perfect matching. (1 mark)